

ON AN ECONOMICAL CONSTRUCTION OF THE TRANSITIVE CLOSURE OF A DIRECTED GRAPH

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Abstract

Full Text

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MATHEMATICS

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ON AN ECONOMICAL CONSTRUCTION OF THE TRANSITIVE CLOSURE OF A DIRECTED GRAPH

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1. For a directed graph (H, γ) , where H is the set of vertices and γ is a mapping of H into itself ⁽¹⁾, algorithms are known for constructing the transitive closure Γ with an estimate of the number of operations, for a graph of general form, $O(n^3)$ ⁽²⁻⁴⁾, where $n = |H|$. In ⁽⁴⁾ the problem of constructing the transitive closure of a graph of general form is reduced to a sequence of three problems: constructing the Hertz graph ⁽⁵⁾ of the given graph, the transitive closure of the Hertz graph (acyclic), and constructing the transitive closure of the given graph from the transitive closure of its Hertz graph; moreover, it is shown that the first and third of these can be solved in $O(n^2)$ operations.

In the present note an algorithm is constructed for the transitive closure of an **acyclic** graph in $O(n^3/\ln n)$ operations.

2. For an acyclic graph (H, γ) , consider the partition of H into ranks K_i :

$$K_0 = \{h \in H : \gamma^{-1}h = \phi\}, \quad K_i = \{h \in H \setminus S_{i-1} : \gamma^{-1}h \subset S_{i-1}\},$$

where $S_i = \bigcup_{j \leq i} K_j$.

In ⁽⁴⁾ an algorithm for such a partition in $O(n^2)$ operations is given. Denote by γ_i and Γ_i the mappings of S_{i-1} into K_i generated by γ and Γ , respectively, and let $G_i = \bigcup \Gamma_j$ be the mapping of S_{i-1} into S_i . Then, obviously:

$$G_0 = \phi, \quad \Gamma_i = \gamma_i G_{i-1} \cup \gamma_i. \quad (1)$$

Thus the construction of the transitive closure has been reduced to a triangular process of obtaining products of mappings.

3. Lemma (M. Kronrod). Let A, B, C be sets, $|A| = p$, $|B| = q$, $|C| = r$, and let α, β be multivalued mappings $\alpha : A \rightarrow B$, $\beta : B \rightarrow C$.

Then the mapping $\beta\alpha : A \rightarrow C$ can be constructed in $O((p+q)qr/\ln q)$ operations.

Proof. Partition B into $[q/\ln q] + 1$ disjoint subsets B_i so that $|B_i| \leq \ln q$. Denote by $\alpha_i : A \rightarrow B_i$, $\beta_i : B_i \rightarrow C$ the mappings generated by α and β , respectively. Then, obviously:

$$\beta\alpha = \bigcup_i \beta_i\alpha_i. \quad (2)$$

Consider the set $M_i = \{m_{is}\}$ of all subsets of B_i and the set $L_i = \{l_{is}\}$ isomorphic to it, where $l_{is} = \bigcup_{b \in m_{is}} \beta_i b$. Numbering the elements of B_i : b_{i0}, b_{i1} , etc., we order the elements of M_i and L_i as follows:

$$s = \sum_k 2^{j_k} \leftrightarrow m_{is} = \bigcup_k b_{ij_k}.$$

It is obvious that $l_{i0} = \phi$, $l_{i2^j} = \beta_i b_{ij}$, while any other l_{is} can be obtained by taking the union of two others with smaller numbers. For example: $l_{is} = l_{ik} \cup l_{i,s-k}$, $s \neq 2^j$, where $k = \max\{2^j : 2^j < s\}$. Since $|L_i| \leq q$, and $|l_{is}| \leq r$, L_i can be constructed in $O(qr)$, and all L_i in $O(q^2r/\ln q)$ operations.

Further, from the completeness of M_i it follows that, for $a \in A$, $\alpha_i a \in M_i$ and, consequently, $\beta_i \alpha_i a \in L_i$. Thus, after all the L_i have been constructed, obtaining $\beta\alpha$ by (2) requires $O(pqr/\ln q)$ operations.

4. Theorem. For an acyclic directed graph (H, γ) with $|H| = n$, the transitive closure Γ can be constructed in $O(n^3/\ln n)$ operations.

Proof. Apply the algorithm of the lemma to the construction of $\gamma_i G_{i-1}$ in (1). Since γ_i acts from S_{i-1} , with $|S_{i-1}| < n$, into K_i , with $|K_i| = n_i$, while G_{i-1} acts from S_{i-2} , with $|S_{i-2}| < n$, into S_{i-1} , the construction of $\gamma_i G_{i-1}$ can be performed in $O(n^2 n_i / \ln n)$ operations. Summing over the ranks, we obtain the desired estimate.

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