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Abstract

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MATHEMATICS

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ON RECURSIVE FUNCTIONS OF LARGE RANGE

(Presented by Academician P. S. Novikov on 31 III 1967)

Scollem in ⁽¹⁾ showed that every general recursive function (g.r.f.) $f(x)$ can be represented in the form

$$f(x) = \psi(\mu y \{ \varphi(y) = x \}) = \psi(\varphi^{-1}(x)), \quad (1)$$

where ψ, φ are suitable primitive recursive functions (p.r.f.). Scollem further writes that, probably, a necessary and sufficient condition for ψ to be able to represent all g.r.f.'s in the form (1) with some p.r.f. $\varphi(x)$ is again that ψ be a function of large range, but that he did not investigate this. However, as E. A. Polyakov has observed, there is no single such function ψ by means of which it would be possible to represent any one-place g.r.f. in the form (1). Nevertheless Scollem's conjecture proved to be valid for a representation of one-place g.r.f.'s somewhat different from (1).

Lemma 1. *For every p.r.f. of large range $R(x)$ there is a p.r.f. of large range $K(x)$ such that $R(x) = R(K(x))$.*

Let, for each $a \in N = \{0, 1, 2, \dots\}$, all solutions of the equation $R(x) = a$ be written in the sequence

$$x_0^{(a)} < x_1^{(a)} < x_2^{(a)} < \dots \quad (*)$$

Then, if $x = x_j^{(a)}$, put $K(x) = x_i^{(a)}$, where $i = \mu_t \{ \text{rest}(j, P_t) = 0 \}$, $\text{rest}(x, y)$ is the remainder on division of x by y , and P_t is the prime number with number t . Since every $x \in N$ belongs to one of the sequences of type (*), $K(x)$ will be an everywhere-defined p.r.f. of large range:

$$K(x) = \mu_{t \leq x} \{ |(L_1(t) + 1) \cdot \text{sg} |R(t) - R(x)| -$$

$$-(\mu_{i \leq L_1(x)} \{ \text{rest}(L_1(x), P_i) = 0 \} + 1) | = 0 \},$$

where $L_1(x)$ is a p.r.f. of large range which, together with $R(x)$, carries out a simple one-to-one enumeration of all pairs of natural numbers (see ⁽²⁾, p. 136).

Theorem 1. *If $R(x)$ is a p.r.f. of large range, then: 1) for every partial recursive function (p.r.f.) $f(x)$ defined at zero there exist a p.r.f. $F(x)$ and an integer a such that*

$$f(x) = R(F^{-1}(x)) + a \cdot \overline{\text{sg}} x; \quad (2)$$

2) for every partial recursive function $f(x)$ which is not defined at zero, there exist a p.r.f. $F(x)$ and a partial recursive function

$$\rho(x) = \begin{cases} \text{undefined,} & \text{if } x = 0, \\ 0, & \text{if } x > 0, \end{cases}$$

such that

$$f(x) = R(F^{-1}(x)) + \rho(x). \quad (3)$$

Proof. Let $f(x)$ be an arbitrary partial recursive function. Then there is a p.r.f. $F_1(x, y)$ such that

$$f(x) = R(\mu y \{F_1(x, y) = 0\})$$

(see ⁽²⁾, p. 137). By Lemma 1, there exists a p.r.f. of large range $K(x)$ such that $R(x) = R(K(x))$. Take a p.r.f. $L_2(x)$, which together with $K(x)$ carries out a simple one-to-one enumeration of all pairs of natural numbers, and the function

$B(x) = \mu y \{F_1'(x, y) = 0\}$ can be written in the following way:

$$B(x) = K \left(\mu t \left\{ L_2(t) \cdot \overline{\text{sg}} F_1(L_2(t), K(t)) \cdot \prod_{i=0}^{K(t)-1} \text{sg} F_1(L_2(t), i) = x \right\} \right) + \\ + B(0) \cdot \overline{\text{sg}} x,$$

i.e.

$$B(x) = K(F^{-1}(x)) + B(0) \cdot \overline{\text{sg}} x,$$

where

$$F(x) = L_2(x) \cdot \overline{\text{sg}} F_1(L_2(x), K(x)) \cdot \prod_{i=0}^{K(x)-1} \text{sg} F_1(L_2(x), i).$$

Then

$$\begin{aligned} f(x) &= R(K(F^{-1}(x)) + B(0) \cdot \overline{\text{sg}} x) = \\ &= R(K(F^{-1}(x))) + R(B(0) \cdot \overline{\text{sg}} x) - R(0), \end{aligned}$$

$$f(x) = R(F^{-1}(x)) + (R(B(0)) - R(0)) \cdot \overline{\text{sg}} x.$$

If $f(0)$ is defined, then from (4) we obtain (2); if $f(0)$ is not defined, then we obtain (3).

It is also obvious that if, for some p.r.f. $R(x)$, all one-place g.r.f.'s are representable in the form (2), then $R(x)$ must be a function of large span.

Remark. Let M be the collection of all p.r.f.'s of large span which differ from one another only at zero. The following assertion is true: there exists such a set M that for every g.r.f. $f(x)$ there are p.r.f.'s $\psi(x) \in M$ and $\varphi(x)$ such that

$$f(x) = \psi(\varphi^{-1}(x)). \quad (5)$$

The representation (5) is a certain strengthening of the result in (1).

Next we consider the algebras $\mathfrak{A}_{\text{o.r.}} = \langle A_{\text{o.r.}}; +, *, -1 \rangle$ and $\mathfrak{A}_{\text{q.r.}} = \langle A_{\text{q.r.}}; +, *, -1 \rangle$ (see (3)). Let

$$M(x) = a_0 + a_1x + a_2x^2 + \dots + a_mx^m, \quad P(x) = b_0 + b_1x + b_2x^2 + \dots + b_nx^n$$

($m, n > 1$, $a_0 \neq 0$) be polynomials with rational coefficients such that

$$(\forall x)[x \in N \rightarrow M(x) \in N \ \& \ P(x) \in N].$$

Define the function $T_P(x)$ as the distance from x to the nearest number on the right of the form $P(t)$.

Theorem 2. The functions $M(x)$ and $T_P(x)$ form a basis of the algebra $\mathfrak{A}_{\text{q.r.}}$.

Proof. From the functions $M(x)$,

$$x \dot{-} y = T_P(T_P^{-1}(x) + y) = x - y \quad \text{for } x \geq y,$$

$$D_a(x) = \mu y \{ ay = x \} = [x/a] \quad \text{for } x = a \cdot t,$$

with the aid of the operations $+$, $*$, -1 one can obtain all linear functions, x^2 , $x \cdot y$, and a polynomial with integer coefficients $P_1(x) = a \cdot P(x)$, where a is a suitable number from N ;

$$\overline{\text{sg}} x = \begin{cases} 1, & x = 0, \\ 0, & x > 0, \end{cases} \quad x \dot{\div} y = \begin{cases} x - y, & x \geq y, \\ 0, & x < y. \end{cases}$$

Let

$$l_1(x) = \mu y \{P_1(y) = a \cdot (x + T_P(x))\},$$

i.e.

$$l_1(x) = P_1^{-1}(x) * a \cdot (x + T_P(x)).$$

We note that

$$(\exists x_0)(\forall x)[x > x_0 \ \& \ P(i) < x \leq P(i+1) \rightarrow l_1(x) = i + 1],$$

i.e. $l_1(x)$ assumes all natural-number values greater than d_0 ($P(d_0) \geq x_0$), and

$$(\exists t)(\forall x)[x > t \rightarrow P(x) \geq x^2].$$

One can verify that in the sequence

$$\langle l_1(0), T_P(0) \rangle, \langle l_1(1), T_P(1) \rangle, \langle l_1(2), T_P(2) \rangle, \dots$$

there will occur any pair of natural numbers of the form (a^2, a) , where $a > t$. Then for $x > t \vee x = 0$:

$$[\sqrt{x}] = T_P(\mu z \{ (l_1(z) + 1) \cdot \text{sg}(T_P^2(z) \dot{\div} l_1(z)) = x + 1 \}). \quad (6)$$

In order for formula (6) to be valid for all x , it is enough to make a small transformation. Now $q(x) = x \dot{\div} [\sqrt{x}]^2$, and the functions $x + 1$ and $q(x)$ form a basis of the algebra $\mathfrak{A}_{\text{q.r.}}$ (see (2), p. 121).

Remark. Let the function $T_{at^n}(x)$ be equal to the distance from x to the nearest number on the right of the form $a \cdot t^n$ ($a > 0$, $n \geq 2$). Then the function $T_{at^n}(x)$ and an arbitrary q.r.f. $f(x)$, where $f(0) = d \neq 0$, form a basis of the algebra $\mathfrak{A}_{\text{q.r.}}$.

The proof of this assertion can be reduced to the proof of Theorem 3, after first obtaining from $f(x)$ and

$$T_{at^n}^n(x) = a \left(\left[\sqrt[n]{\frac{x}{a}} \right] + \text{sg } x \right)^n - x$$

the function $x^{\tilde{n}}$ ($\tilde{n} > 1$).

For p.r.f. an analogous result was obtained by I. A. Lavrov (4).

It is not hard to see that a basis of the algebra $\mathfrak{A}_{o.r.}$ is also a basis of the algebra $\mathfrak{A}_{p.r.}$. The question arises: will every o.r. basis of the algebra $\mathfrak{A}_{p.r.}$ be a basis of the algebra $\mathfrak{A}_{o.r.}$ (by a basis is meant a system of generating elements of the algebra)?

Theorem 3. *There exists an o.r. basis of $\mathfrak{A}_{p.r.}$ which is not a basis of $\mathfrak{A}_{o.r.}$.*

Proof. From the remark to Theorem 3 it follows, for example, that $x + 1$ and $T_{2t^2}(x)$ form a basis of the algebra $\mathfrak{A}_{p.r.}$. But these functions belong to a proper subalgebra of the algebra $\mathfrak{A}_{o.r.}$ (see ⁽⁵⁾) and therefore cannot be a basis of $\mathfrak{A}_{o.r.}$.

Next we construct examples of partial recursive functions of large span, each of which together with $x + 1$ does not form a basis of the algebra $\mathfrak{A}_{p.r.}$; this will give a negative answer to one of A. I. Mal' tsev' s questions.

Let $\Phi_i(x)$ be equal to the distance from x to the nearest, on the left (or on the right), number of the form i^t , $\Phi_i(0) = 0$ ($i = 2, 3, \dots$). Define the function

$$Q_i(x) = \begin{cases} \text{undefined,} & \Phi_i^{-1}(\Phi_i(x)) = x, \\ \Phi_i(x) & \text{otherwise;} \end{cases}$$

$Q_i(x)$ is undefined at those points x at which $\Phi_i(x)$ assumes, for the first time, one of its values. It is clear that $Q_i(x)$ is a p.r.f. of large span, and there is an interval of arbitrarily large length in which it is nowhere defined.

Theorem 4. *The function $x + 1$ and the p.r.f. of large span $Q_i(x)$ are not a basis of $\mathfrak{A}_{p.r.}$.*

Proof. It is enough to note that the functions $x + 1$ and $Q_i(x)$ belong to the following set C , closed under $+$, $*$, and -1 :

$$f(x) \in C \iff \left\{ (\exists a)(\exists b) \left[\begin{array}{l} a, b \text{ rational,} \\ a > 0, \end{array} \quad \& f(x) = \begin{cases} ax + b, & x \in D_f, \\ \text{undefined} \\ \text{otherwise,} \end{cases} \right] \vee \right. \\ \left. \vee (\forall m)(\forall n) \left[m, n \in N \& m > 0 \& f(mx + n) \text{ is undefined in} \right. \right. \\ \left. \left. \text{infinitely many points} \right] \right\};$$

where D_f is the domain of definition of f .

The question of whether every o.r.f. of large span, together with the function $x + 1$, forms a basis of $\mathfrak{A}_{p.r.}$, remains open.

In the paper ⁽⁶⁾ R. Robinson showed that from the functions $x + 1$,

$$\chi(x) = \begin{cases} 1, & x \neq t^2, \\ 0, & x = t^2 \end{cases}$$

with the aid of the operations $|x - y|$, $*$, and i , one can obtain all unary p.r.f. Let us consider the analogous question for o.r.f.

Theorem 5. *The functions 1 and $[\sqrt{x}]$ form a basis of the algebra*

$$\mathfrak{A}'_{o.r.} = \langle A_{o.r.}; |x - y|, *, -1 \rangle.$$

Let us note that if in $\mathfrak{A}_{p.r.}$ the operation $+$ is replaced by the operation $|x - y|$, then the basis is simplified.

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Note: Figure translations are in progress. See original paper for figures.

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