

# ITERATIVE SOLUTION OF PROBLEMS OF LINEAR AND QUADRATIC PROGRAMMING

MATHEMATICS

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**Abstract**

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**MATHEMATICS**

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## ITERATIVE SOLUTION OF PROBLEMS OF LINEAR AND QUADRATIC PROGRAMMING

*(Presented by Academician L. V. Kantorovich, 19 VII 1966)*

In the present note, problems of linear and quadratic programming are solved by a method that is an extension of the method of steepest descent <sup>(1)</sup> to certain extremal problems among whose constraints there are inequalities.

Consider the linear programming problem:  
find the minimum

$$\sum_{j=1}^n c_j x_j$$

under the conditions

$$x_j \geq 0, \quad j = 1, 2, \dots, n; \quad (1)$$

$$\sum_{j=1}^n a_{ij} x_j = b_i, \quad i = 1, 2, \dots, m, \quad (2)$$

where  $c = (c_1, c_2, \dots, c_n)$  and  $x = (x_1, x_2, \dots, x_n)$  are vectors of the  $n$ -dimensional Euclidean space  $E_n$ ; the rank of the matrix  $\|a_{ij}\|$  is equal to  $m$ ;  $b = (b_1, b_2, \dots, b_m)$  is a vector from  $E_m$ . Let  $T \subset E_n$  be the set of vectors satisfying conditions (1) and (2).

We shall call the **dual estimates** of  $x \in T$  the system of numbers

$$u(x) = (u_1(x), u_2(x), \dots, u_m(x)),$$

satisfying the system of equations

$$\sum_{t=1}^m b_{st} u_t = d_s, \quad (3)$$

where

$$b_{st} = \sum_{j=1}^n x_j^2 a_{sj} a_{tj}, \quad d_s = \sum_{j=1}^n x_j^2 c_j a_{sj}, \quad s = 1, 2, \dots, m.$$

**Remark.** The vector  $(u_1(x), u_2(x), \dots, u_m(x))$  is the solution of the problem: find

$$\min_{u \in E^m} \sum_{j=1}^n x_j^2 \left( \sum_{i=1}^m a_{ij} u_i - c_j \right)^2.$$

Let  $U(x)$  be the set of all solutions of system (3). Note that if  $x_j > 0$  for  $j = 1, 2, \dots, n$ , then system (3) has a unique solution.

Set

$$\Phi(x) = \sum_{j=1}^n x_j^2 \left( \sum_{i=1}^m a_{ij} u_i(x) - c_j \right)^2,$$

$$s_j(x) = x_j^2 \left( \sum_{i=1}^m a_{ij} u_i(x) - c_j \right).$$

We now give the algorithm for solving the linear programming problem. Let  $x_j^0 > 0$  for  $j = 1, 2, \dots, n$ ; compute

$$x_j^{k+1} = x_j^k + \lambda_k s_j(x^k)$$

for  $j = 1, 2, \dots, n$ ,

$$\lambda_k = \frac{1}{\sqrt{\Phi(x^k)}}.$$

**Theorem.** If  $\bar{x}$  is a limit point of the sequence  $\{x^k\}$ , then  $\bar{x}$  is a solution of the linear programming problem.

The proof of the theorem follows from the validity of the following assertions, each of which is not difficult to establish.

1.

$$\sum_{j=1}^n a_{ij} s_j(x) = 0, \quad i = 1, 2, \dots, m.$$

2.

$$\sum_{j=1}^n c_{js} j(x) = -\Phi(x).$$

3. If  $x > 0$ , then

$$x + s(x)/\sqrt{\Phi(x)} > 0.$$

4. The mapping  $x \rightarrow U(x)$  is upper semicontinuous.

Suppose a quadratic programming problem is given: minimize

$$\sum_{i=1}^n \sum_{j=1}^n b_{ij} x_i x_j - 2 \sum_{j=1}^n d_j x_j$$

subject to

$$x_j \geq 0, \quad j = 1, 2, \dots, n;$$

$$\sum_{j=1}^n a_{ij} x_j = b_i, \quad i = 1, 2, \dots, m,$$

where the matrix  $\|b_{ij}\|$  is nonnegative definite, the rank of the matrix  $\|a_{ij}\|$  is equal to  $m$ ,  $d = (d_1, d_2, \dots, d_n)$  and  $x = (x_1, x_2, \dots, x_n)$  are vectors in  $E_n$ , and  $b = (b_1, b_2, \dots, b_m)$  is a vector in  $E_m$ .

Set

$$r_j(x) = \sum_{i=1}^n b_{ij} x_i - d_j.$$

Construct a system of equations differing from (3) in that we replace  $c_j$  by  $r_j(x)$ . From the new system find

$$u(x) = (u_1(x), u_2(x), \dots, u_m(x)).$$

Let

$$\Phi(x) = \sum_{j=1}^n x_j^2 \left( \sum_{i=1}^m a_{ij} u_i(x) - r_j(x) \right)^2,$$

$$s_j(x) = x_j^2 \left( \sum_{i=1}^m a_{ij} u_i(x) - r_j(x) \right).$$

The following algorithm is proposed for solving the quadratic programming problem. Let  $x^0 > 0$ ,

$$x^{k+1} = x^k + \lambda_k s(x^k);$$

if

$$\sum_{i=1}^n \sum_{j=1}^n b_{ij} s_i(x^k) s_j(x^k) > 0,$$

compute

$$\lambda'_k = - \frac{\sum_{j=1}^n r_j(x^k) s_j(x^k)}{\sum_{i=1}^n \sum_{j=1}^n b_{ij} s_i(x^k) s_j(x^k)}.$$

If

$$\lambda'_k < 1/\sqrt{\Phi(x^k)},$$

set  $\lambda_k = \lambda'_k$ . In all other cases take

$$\lambda_k = 1/\sqrt{\Phi(x^k)}.$$

A theorem analogous to the one formulated above is valid.

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## CITED LITERATURE

1. L. V. Kantorovich, DAN, **56**, 233 (1947).

*Note: Figure translations are in progress. See original paper for figures.*

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