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# MATHEMATICS

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**Abstract**

**Full Text**

**MATHEMATICS**

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**ON THE DIAMETER OF A STRONGLY CONNECTED GRAPH**

*(Presented by Academician A. D. Aleksandrov, 10 I 1966)*

In this paper finite directed graphs are considered. Following <sup>(1)</sup>, by  $n(G)$ ,  $m(G)$ ,  $\nu(G)$  we denote respectively the number of vertices, the number of arcs, and the cyclomatic number of the graph  $G$ , and at times we shall omit the argument. Obviously,  $m \leq n(n - 1)$ .

Recall that a directed graph  $G$  is called **strongly connected** if for every ordered pair of vertices  $x, y$  there exists a path from  $x$  to  $y$ . For a strongly connected graph  $G$  one always has  $n \leq m$ , and therefore  $1 \leq \nu \leq (n - 1)^2$ .

Let  $\rho(x, y)$  denote the number of arcs of a shortest path leading from  $x$  to  $y$ , and let

$$\rho(G) = \min_x \max_y \rho(x, y)$$

be the radius of the graph  $G$ , and

$$\delta(G) = \max_x \max_y \rho(x, y)$$

the diameter of this graph.

In <sup>(2)</sup> the inequality

$$\rho(G)\nu(G) \geq n(G) - 1. \tag{1}$$

was proved.

Hence follows the estimate\* of the radius

$$\rho(G) \geq \left\lceil \frac{n(G) - 1}{\nu(G)} \right\rceil, \tag{1'}$$

and for every pair of integers  $n, \nu$  ( $1 \leq \nu \leq (n - 1)^2$ ) it is not difficult to give an example of a graph with  $n(G) = n$ ,  $\nu(G) = \nu$ , and

$$\rho(G) = \left\lceil \frac{n - 1}{\nu} \right\rceil.$$

**Theorem.** *For every strongly connected graph  $G$  for which  $\nu(G) \geq 2$ , the relation*

$$\delta(G)\nu(G) \geq 2(n(G) - 1) \tag{2}$$

holds.

From this inequality follows the estimate of the diameter for  $\nu(G) \geq 2$ :

$$\delta(G) \geq \left\lceil \frac{2(n(G) - 1)}{\nu(G)} \right\rceil. \quad (2')$$

Let now  $\mathfrak{G}_{\nu, n}$  be the set of all strongly connected graphs  $G$  with  $n(G) = n$ ,  $\nu(G) = \nu$ , and let

$$\delta(\nu, n) = \min_{G \in \mathfrak{G}_{\nu, n}} \delta(G).$$

Obviously,  $\delta(1, n) = n - 1$  and  $\delta((n - 1)^2, n) = 1$ . If  $n - 1 < \nu < (n - 1)^2$ , then, as is easy to show,  $\delta(\nu, n) = 2$ .

It follows from (2') that for  $2 \leq \nu \leq n - 1$

$$\delta(\nu, n) \geq \left\lceil \frac{2(n - 1)}{\nu} \right\rceil.$$

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\*  $\lceil x \rceil$  denotes the least integer  $a$  for which  $x \leq a$ .

Considering rosettes (see (1), p. 135) with  $n$  vertices and  $\nu$  petals, we obtain

$$\delta(\nu, n) \leq \left\lceil \frac{2(n - 1)}{\nu} \right\rceil + 1,$$

and in the case

$$\nu \geq \frac{n - \delta - 1}{\lceil \delta/2 \rceil},$$

where

$$\delta = \left\lceil \frac{2(n - 1)}{\nu} \right\rceil,$$

we have

$$\delta(\nu, n) = \left\lceil \frac{2(n - 1)}{\nu} \right\rceil.$$

Figure 1 gives an example of a rosette with  $n = 8$  vertices,  $\nu = 3$  petals, and diameter

Fig. 1

Figure 1: Fig. 1

$$\delta(3, 8) = \left\lceil \frac{2(8-1)}{3} \right\rceil = 5.$$

**Fig. 1**

Let us note that the problem of finding an estimate for the diameter of a strongly connected graph using only the numbers of arcs and vertices of the graph, or, equivalently, the number of vertices and the cyclomatic number, was posed in (1). There also a conjecture belonging to Bratton (2) was formulated. Inequality (2'), obtained in the present paper, is stronger than Bratton's conjecture.

Let  $x$  be some vertex of a strongly connected graph  $G$ . Construct a partial graph  $A$  that is a tree growing from the root  $x$ .<sup>\*</sup> We shall always construct the tree  $A$  in such a way that condition  $\mathfrak{A}$  is satisfied: the  $i$ -th level of the tree consists of all vertices  $y$  of the graph  $G$  for which  $\rho(x, y) = i$ . Then the height of the constructed tree is equal to

$$\rho(x) = \max_y \rho(x, y).$$

The following simple result holds.

**Lemma 1.** If  $\lambda$  is the number of terminal vertices of a tree  $A$  satisfying condition  $\mathfrak{A}$ , then

$$\rho(x)\lambda \geq n - 1.$$

Hence, in particular, inequality (1) follows, since always  $\lambda \leq \nu(G)$ .

**Definition.** A path  $\mu$  of a graph  $G$  will be called **straight** if in the graph  $G$  from every vertex of this path, except perhaps the last, exactly one arc leaves.

Let  $s$  be the number of arcs of the longest straight path.

**Lemma 2.** If, for a strongly connected graph  $G$ ,  $\nu(G) \geq 2$  and  $\lambda$  is the number of terminal vertices of a covering tree  $A$  growing from the vertex  $x_0$ , the beginning of some straight path of length  $s$ , then  $\nu \geq \lambda + 1$ .

**Proof.** Since from every vertex of  $G$  at least one arc must leave, we have  $\nu \geq \lambda$ . Suppose  $\nu = \lambda$ . This means that the arcs of the graph  $G$  not belonging to the tree (their number is equal to  $\nu$ ) leave the terminal vertices of the tree, and moreover exactly one arc leaves each terminal vertex. Consequently, at least one of these arcs enters the root vertex  $x_0$ . But this means the existence in the graph  $G$  of a straight path of length  $s + 1$ , which is impossible.

We now turn to the proof of the theorem. For this purpose construct a covering tree  $A$  growing from the vertex  $x_0$ , the initial vertex of some straight path  $\mu_0$  of length  $s$ . As before, let  $\lambda$  denote the number of terminal vertices of the tree  $A$ . Consider two cases.

1.  $s \geq \delta/2$ . Then, as is easy to see,

$$n - 1 \leq s + (\delta - s)\lambda = \delta\lambda - (\lambda - 1)s,$$

and therefore

$$2(n - 1) \leq 2\delta\lambda - 2s(\lambda - 1) \leq 2\delta\lambda - \delta(\lambda - 1) = \delta(\lambda + 1).$$

\* That is, in the terminology of (1), a pradtrees with root  $x$ .

Using Lemma 2, we obtain

$$2(n - 1) \leq \delta\nu.$$

2.  $s < \delta/2$ . In the set  $X$  of all vertices of the tree  $A$ , distinguish three subsets  $X_1, X_2, X_3$ :  $X_1$  is the set of vertices of the path  $\mu_0$ ;  $X_2$  is the set of vertices  $x$  of the tree for which in  $A$  there exists a path of length  $\leq s - 1$  from  $x$  to some terminal vertex of the tree (for  $s = 0$  the set  $X_2$  is empty);  $X_3 = X \setminus (X_1 \cup X_2)$ .

Obviously,

$$n(G) = |X| \leq |X_1| + |X_2| + |X_3|.$$

It is easy to see that

$$|X_1| = s + 1, \quad |X_2| \leq s\lambda.$$

We shall show that  $|X_3| \leq (\delta - 2s)\nu/2$ . Indeed, if  $\mu$  is any path in the tree  $A$  joining two vertices of the subset  $X_3$ , then, obviously, the length of this path is  $\leq \delta - 2s - 1$ , and therefore the number of its vertices is  $\leq \delta - 2s$ . Let  $\sigma$  denote the number of longest paths of the tree  $A$  whose initial and terminal vertices belong to  $X_3$ . It is easy to see that  $\sigma$  is equal to the number of vertices  $z$  in  $X_3$  possessing the following property: every path of the tree  $A$  beginning at the vertex  $z$  contains no other vertices of the subset  $X_3$ . Denote the set of such vertices by  $X'_3$ .

Let  $A(z)$  be the set of vertices of the tree  $A$  reachable in  $A$  from the vertex  $z$ . To each vertex  $z$  we associate the set  $B(z)$  of arcs of the graph  $G$  which do not belong to the tree  $A$  and leave vertices of the set  $A(z)$ . Obviously, if  $z_1 \neq z_2$ , then  $B(z_1) \cap B(z_2) = \emptyset$ .

We shall show that  $|B(z)| \geq 2$  for all  $z \in X'_3$ . Suppose the contrary:  $|B(z_0)| = 1$  for some vertex  $z_0 \in X'_3$ . Then in the tree  $A$  from the vertex  $z_0$  there issues a unique path  $\mu(z_0)$ , ending at a terminal vertex of the tree, and from the vertices of  $\mu(z_0)$  there issues only one arc  $u$  not belonging to the tree. This arc must issue from the terminal vertex of the path  $\mu(z_0)$ . Let  $t$  be the vertex into which the arc  $u$  enters. Then  $t$  belongs to the set of vertices of the path  $\mu(z_0)$ , since otherwise this would mean the existence in the graph  $G$  of a directed path of length  $s + 1$ . But if  $t$  is one of the vertices of the path  $\mu(z_0)$ , then we obtain a circuit  $C$ , from none of whose vertices there issues an arc not belonging to the circuit, which is also impossible, since the graph  $G$  is strongly connected and  $\nu(G) \geq 2$ . Consequently,  $|B(z)| \geq 2$  for all  $z \in X'_3$ , and therefore  $\sigma \leq \nu/2$ .

Thus,

$$n \leq s + 1 + s\lambda + (\delta - 2s)\sigma \leq 1 + s(\lambda + 1) + (\delta - 2s)\nu/2,$$

$$2(n - 1) \leq 2s(1 + \lambda) + (\delta - 2s)\nu \leq 2s\nu + (\delta - 2s)\nu = \delta\nu.$$

The theorem is proved.

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*Note: Figure translations are in progress. See original paper for figures.*

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