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ON INTEGRATION WITH ACCURACY CONTROL

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Abstract

Full Text

CYBERNETICS AND CONTROL THEORY

A. S. KRONROD

ON INTEGRATION WITH ACCURACY CONTROL

(Presented by Academician M. V. Keldysh, 27 VII 1963)

1°. For integrating $f(x)$ on $[-1, +1]$, quadrature formulas of the form

$$A(f) = C_1 f(x_1) + \dots + C_n f(x_n) \quad (x_i \neq x_k \text{ for } i \neq k) \quad (1)$$

are considered. By the **accuracy** of a quadrature A we shall mean the greatest natural number $T(A)$ for which all polynomials of degree $T(A)$ are integrated exactly by formula (1). The number $(A) = n$ will be called the **cost** of A . Among quadratures of cost n , the Gaussian quadrature Γ_n has the maximal accuracy $2n - 1$. The accuracy of other quadratures is $\leq 2n - 2$.

2°. Let a pair of quadratures A, B be given. This means that A and B are different quadratures and $T(A) \leq T(B)$. By the **cost** (A, B) of the pair A, B we shall mean the total number of distinct points x_i occurring in A and B , and by the **accuracy** of the pair the number $T(A, B) = T(A)$.

3°. **Lemma.** Let x_1, \dots, x_k be arbitrary distinct points. There exists a quadrature D

$$D(f) = d_1 f(x_1) + \dots + d_k f(x_k) \quad (2)$$

of accuracy $\geq k - 1$. Among quadratures of accuracy $\geq k - 1$, it is unique.

Indeed, for $f(x) = x^s$ ($s = 0, 1, \dots, k - 1$), the determinant of the system

$$D(x^s) = 2 : (s + 1) \quad (s = 0, 1, \dots, k - 1) \quad (2a)$$

is a Vandermonde determinant, different from zero.

4°. **Theorem 1.** Let A, B be a pair of quadratures. Then

$$(A, B) \geq T(A, B) + 2.$$

Indeed, let $x_s [s = 1, 2, \dots, N = (A, B)]$ be all the points entering into the formulas A and B . The quadratures A and B have the form (2). By the lemma, they are identical if $T(A) \geq (A, B) - 1$.

5°. For integration with accuracy control, with permitted cost $N = 2n + 1^*$, we shall call **optimal** a pair A, B with maximal accuracy $T(A, B)$, minimal cost $(A, B) \leq N$ among all such pairs, with minimal (A) among these, and with maximal $T(B)$ among these. According to 1° and 4° we have $A \equiv \Gamma_n$. Therefore $T(A, B) = 2n - 1 = N - 2$. We construct a quadrature K_n which, together with Γ_n , gives an optimal pair. Let L be the Legendre polynomial of degree n with leading coefficient 1. Let I_p be the $(n + p)$ -th moment of L . Let $t = 0$ for even n , and $t = 1$ for odd n . Put

$$b_0 = 1; \quad b_{2k} = -(b_0 I_{2k} + b_2 I_{2k-2} + \dots + b_{2k-2} I_2) : I_0; \quad (3)$$

$$\varphi(x) = b_0 x^{n+1} + b_2 x^{n-1} + \dots + b_{n+t} x^{1-t}. \quad (4)$$

Then $F(x) = L(x) \cdot \varphi(x)$ is orthogonal to x^s ($s = 0, 1, \dots, n + 1 - t$). The function $F(x)$ has $2n + 1$ distinct roots $\{\beta_s\}$ on $[-1, +1]**$; moreover—

* For an even permitted cost, I do not know how to approach this problem reasonably.

** For arbitrary n I do not know how to prove this, but up to $n \leq 40$ the roots have been computed directly.

No.	Nodes	Gauss quadrature weight	Exact quadrature weight	Relative error in the last significant digit
2	0.0370	0.0000	0.0989	$r_4 = 0.444 \cdot$
	8995	0000	8989	$10^0 r_8 =$
	01142113	00005000	89902454	0.163 ·
	2486	0000	5454	$10^{-1} r_{10} =$
	54055000	00000000	54553111	$0.298 \cdot 10^{-1}$
	0000 0000	0000 0000	1111 1110	
3	0.0197	0.0000	0.0523	$r_6 = 0.160 \cdot$
	5436	0000	2811	$10^0 r_{12} =$
	56461127	00002777	30131342	0.182 ·
	0166	7777	4404	$10^{-2} r_{14} =$
	53792828	77780000	49342006	0.161 ·
	7812	0000	9807	$10^{-1} r_{16} =$
	53275000	00004444	93882254	$0.102 \cdot 10^{-1}$
	0000 0000	4444 4444	8586 9330	

No.	Nodes	Gauss quadrature weight	Exact quadrature weight	Relative error in the last significant digit
4	0.0117	0.0000	0.0314	$r_8 = 0.523 \cdot$
	1287	0000	6813	$10^{-1}r_{14} =$
	46319514	00001739	68330858	0.755 ·
	9308	7391	2680	$10^{-4}r_{16} =$
	45411798	22590000	26661333	0.491 ·
	5689	0000	9917	$10^{-3}r_{18} =$
	12523300	00003260	02261634	0.152 ·
	0994	7257	2459	$10^{-2}r_{20} =$
	82085000	74310000	48011732	$0.330 \cdot 10^{-2}$
	0000 0000	0000 0000	2149 0944	
5	0.0079	0.0000	0.0212	$r_{10} = 0.162 \cdot$
	5731	0000	9101	$10^{-1}r_{18} =$
	99530786	00001184	83760565	0.466 ·
	5805	0634	4165	$10^{-5}r_{20} =$
	70311229	42280000	53110934	0.582 ·
	1663	0000	0039	$10^{-4}r_{22} =$
	67152307	00002393	82781205	0.207 ·
	0534	0438	2016	$10^{-3}r_{24} =$
	49473601	52500000	96141364	0.720 ·
	8979	0000	2482	$10^{-3}r_{26} =$
34195000	00002841	89561441	$0.106 \cdot 10^{-2}$	
0000 0000	4444 4444	9370 8930		
6	0.0055	0.0000	0.0151	$r_{12} = 0.480 \cdot$
	4839	0000	9807	$10^{-2}r_{20} =$
	86940397	00000856	266004418	0.634 ·
	6524	6224	4722	$10^{-6}r_{22} =$
	28980893	61900000	02230686	0.446 ·
	1332	0000	6030	$10^{-5}r_{24} =$
	95671693	00001803	23170905	0.800 ·
	3530	8978	9599	$10^{-4}r_{26} =$
	67672634	65240000	71621106	0.535 ·
	4849	0000	0690	$10^{-4}r_{28} =$
	73623806	000002339	46381168	0.129 ·
	9040	5696	8543	$10^{-3}r_{30} =$
69585000	72860000	20581205	$0.266 \cdot 10^{-3}$	
0000 0000	0000 0000	3629 0088		

No.	Nodes	Gauss quadrature weight	Exact quadrature weight	Relative error in the last significant digit
7	0.0044	0.0000	0.0114	$r_{14} = 0.139 \cdot$
	7231	0000	6766	$10^{-2}r_{16} =$
	44400254	00000647	10050315	0.588 ·
	6604	4248	4604	$10^{-2}r_{24} =$
	38290675	30840000	63150523	0.718 ·
	3787	0000	9500	$10^{-6}r_{26} =$
	83201239	00001398	51610739	0.515 ·
	2493	5263	2066	$10^{-5}r_{28} =$
	22002069	57450000	95980845	0.238 ·
	5638	0000	0236	$10^{-5}r_{30} =$
	22662970	00001502	33200951	0.756 ·
	7742	5253	7528	$10^{-5}r_{32} =$
	43113961	00002089	90321022	0.995 ·
	0725	7959 1836	1617	$10^{-4}r_{34} =$
	24965000		00381047	$0.434 \cdot 10^{-4}$
0000 0000		4107 0542		
8	0.0033	0.0000	0.0089	$r_{16} = 0.396 \cdot$
	1006	0000	1119	$10^{-3}r_{26} =$
	20590198	00000506	16600247	0.645 ·
	5507	1426	1969	$10^{-8}r_{28} =$
	17510529	81450000	75010412	0.539 ·
	3994	0000	4114	$10^{-7}r_{30} =$
	65761016	00001111	94660585	0.242 ·
	0565	1051	2318	$10^{-6}r_{32} =$
	24531658	27270000	54530681	0.842 ·
	6226	0000	8513	$10^{-6}r_{34} =$
	45292372	00001568	46220783	0.229 ·
	3379	5332	2630	$10^{-5}r_{36} =$
	50423196	29390000	30840860	0.540 ·
	4945	0000	3530	$10^{-5}r_{38} =$
	10364082	00001813	42780910	0.110 ·
	8652	4189	2004	$10^{-4}r_{40} =$
	77525000	16890000	23340922	$0.223 \cdot 10^{-4}$
	0000 0000	0000 0000	2320 2872	

No.	Nodes	Gauss quadrature weight	Exact quadrature weight	Relative error in the last significant digit
9	0.0026	0.0000	0.0071	$r_{18} = 0.111 \cdot$
	6091	0000	5238	$10^{-3}r_{30} =$
	96610159	00000406	78220198	$0.738 \cdot$
	1988	3779	1594	$10^{-9}r_{32} =$
	02460425	41810000	75800332	$0.667 \cdot$
	1824	0000	9907	$10^{-8}r_{34} =$
	63750819	00000909	79700475	$0.310 \cdot$
	4841	3040	3955	$10^{-7}r_{36} =$
	43331327	33370000	00450558	$0.118 \cdot$
	5661	0000	9456	$10^{-6}r_{38} =$
	74081933	00001303	73420650	$0.343 \cdot$
	1428	0534	0070	$10^{-6}r_{40} =$
	36502622	82010000	34280726	$0.867 \cdot$
	6883	0000	1979	$10^{-6}r_{42} =$
	44243378	00001567	41920776	$0.191 \cdot$
	7282	7535	8468	$10^{-5}r_{44} =$
	82984178	88380000	39840814	$0.388 \cdot$
	8821	0000	4311	$10^{-5}r_{46} =$
	81935000	00001651	37200824	$0.733 \cdot 10^{-5}$
	0000 0000	1967 7502	4800 6414	

No.	Nodes	Gauss quadrature weight	Exact quadrature weight	Relative error in the last significant digit
10	0.0021	0.0000	0.0058	$r_{20} = 0.307 \cdot$
	7741	0000	4731	$10^{-4}r_{32} =$
	84870130	00000333	94340162	$0.726 \cdot$
	4675	3567	7908	$10^{-10}r_{34} =$
	52410345	21540000	11540272	$0.107 \cdot$
	2473	0000	7987	$10^{-8}r_{36} =$
	21380674	00000747	82270375	$0.374 \cdot$
	6831	2567	1983	$10^{-8}r_{38} =$
	66561095	45750000	74050465	$0.142 \cdot$
	9113	0000	6272	$10^{-7}r_{40} =$
	67071602	00001095	72920541	$0.435 \cdot$
	9251	4318	9357	$10^{-7}r_{42} =$
	58502188	12580000	94010612	$0.114 \cdot$
	1541	0000	4795	$10^{-6}r_{44} =$
	36562833	00001346	81310673	$0.262 \cdot$
	0230	5835	5460	$10^{-6}r_{46} =$
	29953528	96550000	86560713	$0.569 \cdot$
	0358	0000	8796	$10^{-6}r_{48} =$
	65494255	00001477	92890738	$0.112 \cdot$
	6283	6211	8358	$10^{-5}r_{50} =$
	05095000	23580000	74510741	$0.202 \cdot 10^{-5}$
	0000 0000	0000 0000	2227 7000	

No.	Nodes	Gauss quadrature weight	Exact quadrature weight	Relative error in the last significant digit
11	0.0018	0.0000	0.0048	$r_{22} = 0.842 \cdot$
	1519	0000	8272	$10^{-5}r_{24} =$
	30550108	00000278	05230135	$0.525 \cdot$
	8567	3428	7827	$10^{-4}r_{26} =$
	09270291	35580000	73410229	$0.184 \cdot$
	6444	0000	1468	$10^{-3}r_{36} =$
	57710564	00000623	92820325	$0.177 \cdot$
	8680	9018	4871	$10^{-11}r_{38} =$
	41110979	74730000	32750393	$0.881 \cdot$
	7172	0000	2323	$10^{-10}r_{40} =$
	16721349	00000931	96560464	$0.496 \cdot$
	2399	4510	7654	$10^{-9}r_{42} =$
	72131847	94640000	92980529	$0.200 \cdot$
	0023	0000	3603	$10^{-8}r_{44} =$
	99192404	00001165	72410583	$0.610 \cdot$
	4353	3888	2975	$10^{-8}r_{46} =$
	87973010	22260000	71310629	$0.179 \cdot$
	2792	0000	7939	$10^{-7}r_{48} =$
	95243652	00001314	95500656	$0.437 \cdot$
	2842	0227	4034	$10^{-7}r_{50} =$
20244319	22550000	21150672	$0.969 \cdot$	
4390	0000	9678	$10^{-7}r_{52} =$	
96005000	00001364	64000628	$0.199 \cdot 10^{-6}$	
	0000 0000	6254 3390	8889 7350	

whose roots are symmetric with respect to the point $\beta_{n+1} = 0$. Put

$$K_n(f) = d_1 f(\beta_1) + \dots + d_{2n+1} f(\beta_{2n+1}). \quad (5)$$

By the lemma we choose d_i so that $T(K'_n) \geq 2n$. Such a quadrature formula is unique and $d_i = d_{2n+2-i}$. Hence it follows that $T(K_n) \geq 2n + 1$.

6°. Theorem 2. $T(K_n) = 3n + 1$ for even n , and $T(K_n) = 3n + 2$ for odd n .

Indeed, $F(x)$ is orthogonal to x^s ($s = 0, 1, \dots, n + 1 - t$). Further, $\{\beta_i\}$ are the roots of $F(x)$. Hence K_n is exact for certain polynomials of degree $2n + 1 + s$ ($s = 0, 1, \dots, n + 1 - t$) with leading coefficient 1, namely for $F(x) \cdot x^s$. Consequently, K_n is exact for all polynomials up to degree $(3n + 1 + t)$.

n	Nodes	Weights of Gaussian quadrature	Weights of the refining quadrature	Relative error of the refining quadrature
12	0.0015 3303 8735 0092 1903 0287 0027 3497 8624 7374 0413 0188 0075 0178 7672 9230 0269 0480 6889			$0.230 \cdot 10^{-5} K_{26} =$ $0.155 \cdot 10^{-4} K_{28} =$ $0.520 \cdot 10^{-4} K_{38} =$ $0.866 \cdot 10^{-12} K_{40} =$ $0.972 \cdot 10^{-11} K_{42} =$ $0.586 \cdot 10^{-10} K_{44} =$ $0.205 \cdot 10^{-9} K_{46} =$ $0.855 \cdot 10^{-9} K_{48} =$ $0.242 \cdot 10^{-8} K_{50} =$ $0.632 \cdot 10^{-8} K_{52} =$ $0.145 \cdot 10^{-7} K_{54} =$ $0.311 \cdot 10^{-7} K_{56} =$ $0.621 \cdot 10^{-7}$

n	Nodes	Weights of Gaussian quadrature	Weights of the refining quadrature	Relative error of the refining quadrature
13	0.0013	1691	1508	$0.621 \cdot 10^{-6} K_{28} =$ $0.450 \cdot 10^{-5} K_{30} =$ $0.180 \cdot 10^{-4} K_{42} =$ $0.101 \cdot 10^{-12} K_{44} =$ $0.121 \cdot 10^{-11} K_{46} =$ $0.771 \cdot 10^{-11} K_{48} =$ $0.371 \cdot 10^{-10} K_{50} =$ $0.125 \cdot 10^{-9} K_{52} =$ $0.379 \cdot 10^{-9} K_{54} =$ $0.101 \cdot 10^{-8} K_{56} =$ $0.245 \cdot 10^{-8} K_{58} =$ $0.541 \cdot 10^{-8} K_{60} =$ $0.112 \cdot 10^{-7} K_{62} =$ $0.217 \cdot 10^{-7}$

n	Nodes	Weights of Gaussian quadrature	Weights of the refining quadrature	Relative error of the refining quadrature
14	0.0011 3970 31200058 0809905029018326837036309093304066590720657631028012689028			$0.167 \cdot 10^{-6} K_{30} =$ $0.129 \cdot 10^{-5} K_{32} =$ $0.548 \cdot 10^{-5} K_{44} =$ $0.107 \cdot 10^{-13} K_{46} =$ $0.132 \cdot 10^{-12} K_{48} =$ $0.927 \cdot 10^{-12} K_{50} =$ $0.440 \cdot 10^{-11} K_{52} =$ $0.166 \cdot 10^{-10} K_{54} =$ $0.590 \cdot 10^{-10} K_{56} =$ $0.148 \cdot 10^{-9} K_{58} =$ $0.371 \cdot 10^{-9} K_{60} =$ $0.854 \cdot 10^{-9} K_{62} =$ $0.184 \cdot 10^{-8} K_{64} =$ $0.368 \cdot 10^{-8} K_{66} =$ $0.704 \cdot 10^{-8}$

n	Nodes	Weights of Gaussian quadrature	Weights of the refining quadrature	Relative error of the refining quadrature
15	0.0009 9885 6658 0001 0007 0009 0016 26 88 73 19 08 70 06 30 07 80 00 05 14 8 30 23 8 2 8 6 0 7 8 0 6 8			$0.446 \cdot 10^{-7} K_{32} =$ $0.366 \cdot 10^{-6} K_{34} =$ $0.166 \cdot 10^{-5} K_{48} =$ $0.555 \cdot 10^{-14} K_{50} =$ $0.74 \cdot 10^{-13} K_{52} =$ $0.422 \cdot 10^{-12} K_{54} =$ $0.165 \cdot 10^{-11} K_{56} =$ $0.240 \cdot 10^{-11} K_{58} =$ $0.797 \cdot 10^{-11} K_{60} =$ $0.232 \cdot 10^{-10} K_{62} =$ $0.385 \cdot 10^{-10} K_{64} =$ $0.145 \cdot 10^{-9} K_{66} =$ $0.322 \cdot 10^{-9} K_{68} =$ $0.669 \cdot 10^{-9} K_{70} =$ $0.132 \cdot 10^{-8} K_{72} =$ $0.247 \cdot 10^{-8}$

n	Nodes	Weights of Gaussian quadrature	Weights of the refining quadrature	Relative error of the refining quadrature
16	0.008 8036 2927009209953023700011623762243756800642235633010511240716122006710492188			$0.119 \cdot 10^{-7} K_{34} =$ $0.104 \cdot 10^{-6} K_{36} =$ $0.492 \cdot 10^{-6} K_{50} =$ $0.239 \cdot 10^{-15} K_{52} =$ $0.213 \cdot 10^{-14} K_{54} =$ $0.150 \cdot 10^{-13} K_{56} =$ $0.757 \cdot 10^{-13} K_{58} =$ $0.320 \cdot 10^{-12} K_{60} =$ $0.111 \cdot 10^{-11} K_{62} =$ $0.391 \cdot 10^{-11} K_{64} =$ $0.917 \cdot 10^{-11} K_{66} =$ $0.227 \cdot 10^{-10} K_{68} =$ $0.522 \cdot 10^{-10} K_{70} =$ $0.112 \cdot 10^{-9} K_{72} =$ $0.228 \cdot 10^{-9} K_{74} =$ $0.442 \cdot 10^{-9} K_{76} =$ $0.816 \cdot 10^{-9}$

7°. In the table, the second column gives the nodes, i.e. the roots of $F(x)$, for $2 \leq n \leq 16$. In the third and fourth columns the weights of the Gaussian and refining quadratures are given. All pertain to the case of integration on $[0, 1]$. The nodes and weights, by virtue of their symmetry with respect to 0.5, are given only for the left half of the interval $[0, 1]$.

In the last column of the table are given the relative errors $\Gamma_p^{(n)}$ and $K_p^{(n)}$ of the Gaussian and refining quadratures in computing the integral of x^p on $[-1, +1]$ for the first even p for which the quadratures Γ_n and K_n are no longer exact.

All computations were carried out on the machine of the Institute of Theoretical and Experimental Physics. Programs from the library of long-range [[unclear: continuation on next page]] were used.

of A. V. Uskov' s floating-point library and A. Zhivotovsky and V. Pruss' s integer library.

Remark 1. A pair of quadrature formulas—a Gaussian one and a refining one—is convenient to use in solving integral equations with accuracy control in the metric C , since the even nodes of the refining quadrature formula are the nodes of the Gaussian quadrature formula.

Remark 2. A comparison of the corresponding remainder terms $\Gamma_p^{(2n+1)}$ and $K_p^{(n)}$ of the Gaussian and refining quadrature formulas (for example, $\Gamma_{30}^{(15)}$, $\Gamma_{32}^{(15)}$, $\Gamma_{34}^{(15)}$ and $K_{30}^{(7)}$, $K_{32}^{(7)}$, $K_{34}^{(7)}$) shows that, if the required accuracy of integration is not very high, then the refining quadrature formula is only slightly inferior to the Gaussian quadrature formula of the same cost.

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Note: Figure translations are in progress. See original paper for figures.

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