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Abstract

Full Text

MATHEMATICS

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ESTIMATES OF THE NUMBER OF STATES ARISING IN THE DETERMINIZATION OF A NONDETERMINISTIC AUTONOMOUS AUTOMATON

(Presented by Academician A. N. Kolmogorov on 14 XI 1963)

Let \mathfrak{A} be a nondeterministic autonomous automaton with states $\sigma_1, \sigma_2, \dots, \sigma_n$, i.e., a multivalued mapping of the set S of states into itself*:

$$\sigma_i \rightarrow S_i; \quad S_i \subset S \quad (i = 1, 2, \dots, n). \quad (1)$$

We fix one of the states, for example σ_1 , as the initial state.

The mapping (1) is directly extended to the system Σ of all subsets of the set S :

$$M \rightarrow M' = \bigcup_{\sigma_i \in M} S_i \quad (M \in \Sigma). \quad (2)$$

Restricting the mapping (2) to the smallest subsystem $\Sigma_0 \subset \Sigma$ invariant with respect to it and containing the "initial" set $\{\sigma_1\}$, we obtain a certain already deterministic autonomous automaton \mathfrak{A}_0 . The passage from \mathfrak{A} to \mathfrak{A}_0 is naturally called the determinization of the (initial) automaton \mathfrak{A} .

The determinization procedure is used, for example, in ⁽¹⁾ in constructing an automaton representing a given regular event**. In this and in certain other related problems there arises the question of estimating the number of states of the automaton \mathfrak{A}_0 . We shall denote this number by $d(\mathfrak{A})$. We shall obtain for $d(\mathfrak{A})$ an estimate substantially better than the trivial estimate $d(\mathfrak{A}) \leq 2^n$. Moreover, we shall show how $d(\mathfrak{A})$ can be estimated while taking into account the structure of the automaton \mathfrak{A} .

Consider the Boolean algebra of square matrices of order n with entries 0, 1, and associate with the automaton \mathfrak{A} the matrix*** $A = (a_{ik})_{i,k=1}^n$, putting $a_{ik} = 1$ if $\sigma_i \in S_k$, and $a_{ik} = 0$ otherwise. Next consider the sequence of powers

$$A^m = (a_{ik}^{(m)})_{i,k=1}^n \quad (m = 0, 1, 2, \dots). \quad (3)$$

It is easy to see that the states of the automaton \mathfrak{A}_0 are all possible sets of the form

$$\bigcup_{i=1}^n a_{i1}^{(m)} \{\sigma_i\} \quad (m = 0, 1, 2, \dots).$$

Here we put $\varepsilon M = M$ for $\varepsilon = 1$ and $\varepsilon M = \emptyset$ for $\varepsilon = 0$ for all $M \in \Sigma$. It is now clear that

$$d(\mathfrak{A}) \leq l(A), \quad (4)$$

where $l(A) = t(A) + p(A)$, $p(A)$ is the length of the least period of the sequence (3), and $t(A)$ is the length of its preperiod.

* That is, a finite graph.

** Our case of an autonomous automaton corresponds to the representation of events over a one-letter alphabet.

*** The adjacency matrix in the terminology of graph theory.

To estimate the value $l(A)$ we use the theory of matrices with nonnegative entries (²⁻⁶), a number of whose concepts and theorems carry over directly to the Boolean matrices under consideration, and consequently also to nondeterministic autonomous automata.

From one estimate of H. Wielandt (³)*, with the aid of (4), it follows directly that

Theorem 1. *If \mathfrak{A} is a primitive automaton, then*

$$d(\mathfrak{A}) \leq n^2 - 2n + 3. \quad (5)$$

This interpretation of Wielandt' s estimate is essentially known (^{7,8}).

Relying on Wielandt' s estimate, one can also prove the following assertion.

Theorem 2. *If \mathfrak{A} is an imprimitive automaton with index of imprimitivity h , then*

$$d(\mathfrak{A}) \leq \frac{1}{h}(n^2 - 2nh + 4h^2). \quad (6)$$

Comparison of the estimates (5), (6) leads to the following conclusion.

Corollary. *For any irreducible automaton \mathfrak{A} with $n \geq 5$, estimate (5) holds.*

Estimate (5) is sharp in the class of irreducible (even primitive) automata ⁽³⁾.

In the case of a reducible automaton \mathfrak{A} , the estimate for $d(\mathfrak{A})$ becomes much worse in comparison with estimate (5).

Lemma. *If the matrix A has the form*

$$A = \begin{pmatrix} U & 0 \\ W & V \end{pmatrix},$$

where U, V are square blocks, then

$$t(A) < t(U) + l(V) + m(p(U), p(V))$$

and $p(A) = m(p(U), p(V))$, where $m(\ , \)$ denotes the least common multiple.

On the basis of this lemma one establishes

Theorem 3. *For any automaton \mathfrak{A} ,*

$$d(\mathfrak{A}) < m(h_1, h_2, \dots, h_r) + \sum_{k=2}^r m(h_1, h_2, \dots, h_k) + \sum_{k=1}^r \frac{1}{h_k} (v_k^2 - 2v_k h_k + 4h_k^2), \quad (7)$$

where v_1, v_2, \dots, v_r are the numbers of states of the irreducible components of the automaton \mathfrak{A} , and h_1, h_2, \dots, h_r are their indices of imprimitivity.

Of the numerous consequences of Theorem 3 we note the three most interesting.

Corollary 1. *If all irreducible components of the automaton \mathfrak{A} are primitive, then for $n \geq 4$ estimate (5) holds.*

Corollary 2. *For any automaton \mathfrak{A} ,*

$$d(\mathfrak{A}) < r \left(\frac{n}{r} \right)^r + (n - r)^2 + 5r, \quad (8)$$

where r is the number of irreducible components of the automaton \mathfrak{A} .

Corollary 3. *For any automaton \mathfrak{A} with $n \geq 6$, the estimate holds*

$$d(\mathfrak{A}) < M(n) + n^2 - 2n + 6, \quad (9)$$

* Which is a refinement of an earlier estimate of H. Frobenius ⁽²⁾. Wielandt states his estimate without proof

where $M(n)$ is an arithmetic function defined as the maximum of the expression

$$m(N_1, N_2, \dots, N_s) + \sum_{k=2}^s m(N_1, N_2, \dots, N_k)$$

over all systems (N_1, N_2, \dots, N_s) ($s = 2, 3, \dots$) of natural numbers satisfying the condition

$$\sum_{k=1}^s N_k \leq n. \quad (10)$$

An estimate close to (9) was found earlier (but not published) by V. G. Bodnar-chuk and A. A. Letichevskii, on the basis of other considerations.

It is not difficult to prove that, although the function $M(n)$ grows faster than any power n^β , it nevertheless grows more slowly than any exponential $e^{\varepsilon n}$ ($\varepsilon > 0$). These facts were refined by I. V. Ostrovskii, who proved that

$$\ln M(n) \sim \sqrt{n \ln n} \quad (n \rightarrow \infty). \quad (11)$$

Thus, $M(n)$ is the principal term of estimate (9).

Estimate (9) is “almost” exact. Namely, for every n there exists an automaton $\mathfrak{A}^{(0)}$ with n states such that $d(\mathfrak{A}^{(0)}) \geq M_1(n)$, where $M_1(n)$ is an arithmetic function such that $\ln M_1(n) \sim \sqrt{n \ln n}$.

Let us single out some classes of decomposable, generally speaking, automata for which estimate (9) can be substantially improved.

Suppose the number r of components of the automaton \mathfrak{A} is not too large, namely, satisfies the inequality

$$r \leq q \sqrt{\frac{n}{\ln n}} \quad (q = \text{const}, q < 2). \quad (12)$$

Then, by virtue of (8) and (11),

$$d(\mathfrak{A}) \leq o(M(n)). \quad (13)$$

The same result is obtained if r is not too small, namely,

$$r \geq n - q \sqrt{n \ln n} \quad (q = \text{const}, q < 1). \quad (14)$$

Strengthening requirement (14) to

$$r \geq n - \alpha \ln n \quad (\alpha = \text{const}), \quad (15)$$

we obtain the power estimate

$$d(\mathfrak{A}) \leq O(n^\beta), \quad (16)$$

where $\beta = \max(\alpha + 1, 2)$. Estimate (16) is also valid for $r \leq \beta$. In particular, if $r \geq n - \ln n$ or $r \leq 2$, then

$$d(\mathfrak{A}) \leq O(n^2). \quad (17)$$

In conclusion, let us note that, from a methodological point of view, the present work is related to the investigations ^(9, 10).

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