



---

Soviet-era science, translated into English

# MATHEMATICS

1964

SovietRxiv

---

View the original and related papers at <https://sovietrxiv.org/items/ru-196401.38824>

Source: Math-Net.Ru and CyberLeninka. Machine translation. Verify with the original.

Abstract

Full Text

## MATHEMATICS

G. P. Akilov, A. M. Rubinov

### A METHOD OF SUCCESSIVE APPROXIMATIONS FOR FINDING THE POLYNOMIAL OF BEST APPROXIMATION

(Presented by Academician V. I. Smirnov on 26 II 1964)

1. Let  $X$  be a normed space (over the field of all real or all complex numbers); let  $x_1, x_2, \dots, x_n$  be linearly independent vectors in  $X$ ; denote by  $\Omega = \mathcal{L}(\{x_k\})$  the elements of  $\Omega$ —linear combinations of the vectors  $x_1, x_2, \dots, x_n$ —which we shall call **polynomials**; if  $u$  is a given element of the space  $X$ , then a polynomial  $Q$  such that, for  $P = Q$ , the functional  $\|u - P\|$  ( $P \in \Omega$ ) attains its least value  $\mu = \rho(u, \Omega)$ , is called the **polynomial of best approximation** (to  $u$ ), and the number  $\mu$  the **best approximation**.
2. The process of determining the polynomial of best approximation consists in the successive solution of a number of auxiliary problems: let  $g_k \in X^*$  ( $k = 1, 2, \dots, m$ ;  $m \geq n$ ); it is required to determine  $P' \in \Omega$  such that

$$\max_{k \leq m} |g_k(P' - u)| = \min_{P \in \Omega} \left[ \max_{k \leq m} |g_k(P - u)| \right]. \quad (1)$$

As is known, such problems can be solved by means of linear programming, for example with the aid of a certain modification of the simplex method.

3. The construction of the polynomial of best approximation  $Q$  and the determination of the best approximation  $\mu$  are carried out according to the following scheme.

Denote by  $X_0$  the subspace of the space  $X$  spanned by the set  $\Omega$  and the element  $u$ , and let  $\Gamma \subset X^*$  be a set of normalized functionals having the property that for every  $x \in X_0$  one can specify such an  $f_0 \in \Gamma$  that  $|f_0(x)| = \|x\|$ .

As is easy to see, in  $\Gamma$  there exist elements  $f_1, f_2, \dots, f_n$  such that the determinant  $|f_i(x_k)|_{i,k=1,2,\dots,n} \neq 0$ , and then, evidently, there exists an “interpolation” polynomial  $P_n \in \Omega$ , determined by the equalities

$$f_i(u - P_n) = 0 \quad (i = 1, 2, \dots, n).$$

It is clear that  $P_n$  is a solution of problem (1), if in it one takes  $m = n$ ,  $g_k = f_k$  ( $k = 1, 2, \dots, n$ ).

Let, for  $m \geq n$ , functionals  $f_1, f_2, \dots, f_m \in \Gamma$  and polynomials  $P_n, P_{n+1}, \dots, P_m \in \Omega$  be defined so that

$$\max_{k \leq s} |f_k(P_s - u)| = \min_{P \in \Omega} \left[ \max_{k \leq s} |f_k(P - u)| \right] \quad (s = n, n + 1, \dots, m), \quad (2)$$

$$|f_{k+1}(P_k - u)| = \|P_k - u\| \quad (k = n, n + 1, \dots, m - 1). \quad (3)$$

The functional  $f_{m+1}$  is found so that (3) is satisfied for  $k = m$ , and the polynomial  $P_{m+1}$  is determined as the solution of problem (1), where  $g_k = f_k$  ( $k = 1, 2, \dots, m + 1$ ).

As a result of the process described, two sequences are obtained:  $\{f_k\}$  ( $f_k \in \Gamma$ ;  $k = 1, 2, \dots$ ) and  $\{P_s\}$  ( $P_s \in \Omega$ ;  $s = n, n + 1, \dots$ ) such that (2) and (3) are satisfied for all  $k, s \geq n$ . Denote  $\mu_m = \|P_m - u\|$  ( $m \geq n$ ).

**Theorem.** The sequence  $\{\mu_m\}$  converges to  $\mu$ . The sequence  $\{P_m\}$  is bounded (in the space  $X$ ). If  $\{P_{m_i}\}$  is any convergent subsequence of the sequence  $\{P_m\}$ , then  $Q = \lim_{i \rightarrow \infty} P_{m_i}$  is a polynomial of best approximation. If the polynomial of best approximation  $Q$  is unique, then the sequence  $\{P_m\}$  itself converges to  $Q$ .

4. Before proving the theorem we give some preliminary considerations.

In the space  $X$  introduce seminorms  $\|\cdot\|_s$  ( $s \geq n$ ), putting

$$\|x\|_s = \max_{k \leq s} |f_k(x)|.$$

Clearly,

$$\|x\|_s \leq \|x\|_{s+1} \leq \|x\| \quad (x \in X; s = n, n + 1, \dots). \quad (4)$$

**Lemma 1.** If  $s > m$ , then  $\mu_m = \|P_m^* - u\|_s$ .

Next, denote

$$\lambda_m = \|P_m - u\|_m \quad (m \geq n).$$

**Lemma 2.**  $\lambda_n \leq \lambda_{n+1} \leq \dots$ .

**Lemma 3.**  $\lambda_m \leq \mu \leq \mu_m$  ( $m \geq n$ ).

**Remark.** If it turns out that  $\lambda_m = \mu_m$  for some  $m$ , then  $\mu = \lambda_m = \mu_m$  and  $P_m$  will be a polynomial of best approximation. In this case the polynomials  $P_k$  ( $k \geq m$ ) will also be polynomials of best approximation.

5. **Proof of the theorem.** It is not hard to see that in the space  $X_0$  the seminorm  $\|\cdot\|_{n+1}$  will be a norm. But then, since  $X_0$  is finite-dimensional, there is a number  $K$  such that  $\|x\| \leq K\|x\|_{n+1}$  ( $x \in X_0$ ); then, on the basis of Lemma 3 and (4),

$$\mu_m = \|P_m - u\| \leq K\|P_m^* - u\|_{n+1} \leq K\|P_m - u\|_m = K\lambda_m \leq K\mu \quad (m \geq n+1),$$

so that the sequence  $\{\mu_m\}$ , and hence also the sequence  $\{P_m\}$ , is bounded. Consider an arbitrary subsequence  $\{\mu_{m_i}\}$  and the corresponding sequence  $\{P_{m_i}\}$ . Passing, if necessary, to a subsequence of this sequence, we may assume that there exists

$$Q = \lim_{i \rightarrow \infty} P_{m_i}.$$

Then

$$\mu_{m_i} = \|P_{m_i} - u\| \rightarrow \|Q - u\|.$$

Further,

$$\lambda_{m_i} \rightarrow \lambda = \lim_{m \rightarrow \infty} \lambda_m.$$

According to Lemma 1, taking (4) into account, we may write

$$\begin{aligned} \mu_{m_i} &= \|P_{m_i} - u\|_{m_i+1} \leq \|P_{m_{i+1}} - u\|_{m_i+1} + \|P_{m_{i+1}} - P_{m_i}\|_{m_i+1} \\ &\leq \lambda_{m_{i+1}} + \|P_{m_{i+1}} - P_{m_i}\| \quad (i = 1, 2, \dots). \end{aligned}$$

Passing here to the limit, we obtain

$$\mu \leq \|Q - u\| \leq \lambda.$$

But, by Lemma 3,  $\lambda \leq \mu$ . Thus  $\|Q - u\| = \mu$ ,  $\mu_{m_i} \rightarrow \mu$ . It follows that  $\mu_m \rightarrow \mu$ .

The remaining assertions of the theorem follow in an obvious way from what was said above.

## 6. Remarks.

1°. Since  $\lambda_m \rightarrow \lambda = \mu$ , the result of Lemma 3 can be used to estimate the rate of convergence  $\mu_m \rightarrow \mu$  (and, of course,  $\lambda_m \rightarrow \mu$ ).

2°. The method described above for finding a polynomial of best approximation can also be used in the case when it is required that the coefficients  $a_1, a_2, \dots, a_n$  of the desired polynomial satisfy certain linear restrictions:

$$\sum_{k=1}^n c_i^k a_k \leq C_i \quad (i = 1, 2, \dots, \nu). \quad (5)$$

In this case it is only necessary to replace the set  $\Omega$  by the set of all such polynomials whose coefficients satisfy conditions analogous to (5).

3°. If the process described above is modified somewhat, then in the role of  $\Gamma$  one may consider such a set of normalized functionals that, for each  $x \in X_0$ , one has

$$\|x\| = \sup_{f \in \Gamma} |f(x)|.$$

4°. Let us indicate some particular cases.

- a)  $X = C(T)$ , where  $T$  is a compact topological space. As  $\Gamma$  here one may take the set  $T$ , i.e., more precisely, the set of all positive measures with singleton supports.
- b) If  $T$  is a closed bounded set in Euclidean space and  $X = C^{(r)}(T)$ , then one may take  $\Gamma = T \times \{0, 1, \dots, r\}$ .
- c) If  $X = L^p$  ( $1 \leq p \leq \infty$ ), then in the case  $p < \infty$  one should take for  $\Gamma$  the set of all normalized functionals. For  $p = \infty$ , it is expedient to understand  $\Gamma$  as the set of all absolutely continuous positive normalized measures.

Leningrad State University  
named after A. A. Zhdanov

Received  
30 I 1964

*Note: Figure translations are in progress. See original paper for figures.*

*Source: Math-Net.Ru and CyberLeninka. Machine translation. Verify with the original.*