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Abstract

Full Text

MATHEMATICS

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ON A GENERAL APPROXIMATE METHOD FOR SOLVING LINEAR PROBLEMS

(Presented by Academician I. N. Vekua on 3 November 1961)

Let E_1 and E_2 be linear normed spaces, and let a linear operator A be given with domain $G_1 \subseteq E_1$ and range $G_2 \subseteq E_2$. It is required to construct a sequence $\{x_n\}$ for which the relation

$$\lim_{n \rightarrow \infty} \|y - Ax_n\| = \inf_{x \in G_1} \|y - Ax\|, \quad y \in E_2. \quad (1)$$

is satisfied.

Suppose that in G_1 there exists an A -complete system e_1, e_2, \dots, e_n ⁽¹⁾, and we shall seek x_n in the form

$$x_n = \sum_1^n \alpha_k e_k, \quad (2)$$

(α_k may always be assumed real) from the condition that the sought values of the parameters α_k^0 minimize

$$\Delta(P) = \left\| y - \sum_1^n \alpha_k y_k \right\|, \quad y_k = Ae_k \neq 0, \quad P(\alpha_1, \alpha_2, \dots, \alpha_n). \quad (3)$$

It is not difficult to show that

$$\lim_{n \rightarrow \infty} \left\| y - \sum_1^n \alpha_k^0 y_k \right\| = \inf_{x \in G_1} \|y - Ax\|. \quad (4)$$

Further, it is easy to see that

$$\Delta \left[\frac{1}{2}(P_1 + P_2) \right] \leq \frac{1}{2}[\Delta(P_1) + \Delta(P_2)],$$

i.e. Δ is a convex function. Therefore it is natural, for finding α_k^0 , to apply various methods of steepest descent.

We consider an approximate method X , consisting of three principal parts.

1. **Determination of $\inf \Delta$ on a given interval $[P', P'']$.**

Put $\Delta(t) = \Delta(P)$, $P = P' + t(P'' - P')$. Having computed Δ at three values $t = t^{(1)} = \frac{1}{2}$, $t^{(1)} \pm \frac{1}{2}$, choose $t = t^{(2)}$ for which Δ is smallest. If all three values are equal, then they coincide with the required one. Otherwise, repeating the indicated operation for the values $t = t^{(2)}$, $t^{(2)} \pm \frac{1}{2}$, we find $t = t^{(3)}$, for which Δ is smallest, and so on.*

2. **Determination of $\inf \Delta$ in a given direction $-\nabla$.**

Let P_1 be the initial point. Form the sequence of points

$$T_r = P_1 + 2^{r-2}\sigma\nabla, \quad r = 2, 3, \dots \quad (T_0 = T_1 = P_1),$$

where σ is a given positive number, and compute

$$\gamma_r = \Delta(T_r) - \Delta(T_{r-1}), \quad r = 2, 3, \dots$$

Let r_0 , $r_0 \geq 2$, be the first value for which $\gamma_{r_0} \geq 0$. Then the required value of Δ is attained on the interval $[T_{r_0-2}, T_{r_0}]$, and it can be found by algorithm 1.

* If $t^{(j)} > 1$ or $t^{(j)} < 0$, then the required value is attained, respectively, at $t = 1$ or $t = 0$.

Remark 1. For functions Δ of the form (3), Problem 2 reduces to determining the least value of the function $\varphi(a) = \|u - av\|$, where $v \neq \theta$ and u are given elements. It is not difficult to conclude that in this case the desired value is attained on the segment $[-2\|u\|/\|v\|, 2\|u\|/\|v\|]$, and Algorithm 1 determines it in the number of steps $k = \log_2 4\|u\|/h$ with an accuracy not exceeding h .

3. **Determination of the direction of decrease of Δ .** Suppose that there is an algorithm for searching through points which determines $\inf \Delta$ along the boundary Γ_n of an n -dimensional cube with center at the given point P' , with faces parallel to the coordinate planes and with edge length 2δ . As soon as, in the course of solving this problem, a point $P'' \in \Gamma_n$ is obtained for which $\Delta(P') - \Delta(P'') > \sigma$, we shall regard the ray

$$P = P' + t(P'' - P'), \quad t \geq 0,$$

as the desired direction.

Noting that for any points P_1 and P_2

$$|\Delta(P_1) - \Delta(P_2)| \leq \rho(P_1, P_2) \left(\sum_1^n \|y_k\|^2 \right)^{1/2} \leq C_1 \rho(P_1, P_2),$$

and denoting by P'^0 the point nearest to P' at which $\Delta(P'^0) = \inf \Delta$ over all $\{\alpha_k\}$, let us join P' with P'^0 by a straight-line segment. Then, if

$$\Delta(P') - \Delta(P'^0) > C_1 \delta \sqrt{n},$$

then along this segment Δ will decrease, and, by the convexity of Δ ,

$$\Delta(P') - \Delta(P'') \geq \frac{\delta}{\rho(P', P'')} [\Delta(P') - \Delta(P'^0)], \quad (5)$$

where P'' is the point of intersection of the segment with Γ_n .

Let $P_1 = P'$ and $P_2 = P' + t^0(P'' - P')$, where

$$\Delta[P' + t^0(P'' - P')] = \inf_t \Delta[P' + t(P'' - P')].$$

Taking P_2 as the initial point instead of P_1 , in an analogous way we obtain P_3 , then P_4 , etc. Using (5) and the fact that

$$\Delta(P) = \Delta(P^*) = \left\| y - \sum_1^m a_i y_{k_i} \right\|, \quad (6)$$

where $\{y_{k_i}\}$ is a linearly independent system equivalent to $\{y_k\}$, it is not difficult to prove the following theorem:

Theorem 1. *Method X always converges in the sense that, whatever $\varepsilon > 0$, for $\delta < \varepsilon/C_1\sqrt{n}$ and for sufficiently small σ*

$$\Delta(P_j) - \inf_{\{\alpha_k\}} \Delta(P) < \varepsilon \quad (7)$$

for all j , starting from some one.

Remark 2. In some problems (for example, in solving a system of inequalities) the corresponding function Δ , while possessing the property of convexity, does not, generally speaking, possess all the properties of a function of the form (3). Method X can also be extended to these problems, with the same convergence theorem for the method, under the condition that the range of values of all those P for which $\Delta(P) \leq \text{const} = \Delta(P_1)$ is bounded.

The most laborious task is that of determining $\inf \Delta$ on Γ_n . To solve it we choose a sequence of numbers $\sigma \ll \delta_{n-1} \ll \delta_{n-2} \ll \dots \ll \delta_1 = \delta$ and determine the least values of $2n$ convex functions of $n - 1$ variables:

$$\begin{aligned} &\Delta(\alpha_1 \pm \delta, \alpha_2, \dots, \alpha_n), \quad \Delta(\alpha_1, \alpha_2 \pm \delta, \alpha_3, \dots, \alpha_n), \dots, \\ &\dots, \Delta(\alpha_1, \alpha_2, \dots, \alpha_{n-1}, \alpha_n \pm \delta). \end{aligned}$$

We do this by Method X, taking the corresponding $(n - 1)$ -dimensional cubes with edge length $2\delta_2 \ll 2\delta_1$. In this way, by induction, the whole matter is reduced to determining least values in the given directions. As is clear, to implement such an algorithm one needs a computer with a large word length.

For comparatively large n , it will in practice be more advantageous to solve this problem by computing Δ at the nodes of uniform cubic grids with gradually decreasing step size, superimposed on Γ_n . In all those cases when P'^0 lies

outside Γ_n , the largest step h^* for which points P'' having the property $\Delta(P') - \Delta(P'') > \sigma$ are found on the grids may naturally be called a δ, σ -measure of the conditioning of the original problem. The smaller h^* , the worse conditioned the problem is and the more difficult it is to solve.

As an example, consider the problem of determining the optimal parameters of a linear differential equation with constant coefficients of order not higher than the 3rd from a given transition function $h(t)$. Introduce the function

$$\Delta(P) = \max_t \left| \frac{d^3 h}{dt^3} + a_1 \frac{d^2 h}{dt^2} + a_2 \frac{dh}{dt} + a_3 h + a_4 \right|, \quad P(a_1, a_2, a_3, a_4). \quad (8)$$

The problem is to find P^0 for which $\Delta(P)$ attains the smallest value. The function Δ is convex, and by algorithm X the problem can be solved effectively, since the number of undetermined parameters is small. The known algorithms of best approximation ⁽²⁾ are, generally speaking, not applicable in this case, since the system of functions $d^2 h/dt^2$, dh/dt , h , 1 need not be a Chebyshev system.

Let $\Delta(P)$ be an arbitrary everywhere continuously differentiable convex function of n real variables. In seeking $\inf \Delta$ in this case, algorithm 3 can be substantially simplified by taking as the sought directions $-\text{grad } \Delta$ (the method of steepest descent) or, successively, the vectors $e_1(1, 0, \dots, 0)$, $e_2(0, 1, 0, \dots, 0)$, \dots , $e_n(0, \dots, 0, 1)$ (the method of coordinate descent).

Theorem 2. *The method of steepest descent always converges in the sense that*

$$\lim_{j \rightarrow \infty} \Delta(P_j) = \inf_{\{P\}} \Delta(P), \quad (9)$$

where

$$P_j = P_{j-1} + t_j \text{grad } \Delta(P_{j-1}), \quad \Delta(P_j) = \inf_{\{t\}} \Delta[P_{j-1} + t \text{grad } \Delta(P_{j-1})]. \quad (10)$$

Proof. Assuming that $\text{grad } \Delta(P_j) \neq 0$, $j = 1, 2, \dots$, select from $\{P_j\}$ a convergent subsequence $\{P_{j_k}\}$, $\lim_{k \rightarrow \infty} P_{j_k} = \tilde{P}$.

If $\text{grad } \Delta(\tilde{P}) \neq 0$, then

$$\begin{aligned} \Delta(\tilde{P}) &\leq \inf_{\{t_k\}} \Delta[P_{j_k} + t_{j_k+1} \text{grad } \Delta(P_{j_k})] \leq \\ &\leq \lim_{k \rightarrow \infty} \Delta[P_{j_k} + t^0 \text{grad } \Delta(P_{j_k})] = \inf_{\{t\}} \Delta[\tilde{P} + t \text{grad } \Delta(\tilde{P})] < \Delta(\tilde{P}). \end{aligned}$$

Consequently, $\text{grad } \Delta(\tilde{P}) = 0$, and the theorem may be regarded as proved.

Suppose that at every point, except possibly at points where $\Delta(P)$ assumes its smallest value, the minimum of Δ in each of the directions e_s , $s = 1, 2, \dots, n$, is attained only for a single value of the argument. Then the following is true:

Theorem 3. *The method of coordinate descent always converges in the sense that*

$$\lim_{j \rightarrow \infty} \Delta(P_j^s) = \inf_{\{P\}} \Delta(P), \quad s = 1, 2, \dots, n, \quad (11)$$

where

$$P_j^{s+1} = P_j^s + t_j^s e_{s+1}, \quad \Delta(P_j^{s+1}) = \inf_{\{t\}} \Delta(P_j^s + t e_{s+1}), \quad P_j^{n+1} = P_{j+1}^1, \quad e_{n+1} = e_1. \quad (12)$$

Proof. Let $\lim_{k \rightarrow \infty} P_{jk}^s = \tilde{P}^s$, $1 \leq s \leq n$, and let

$$\gamma_i = \inf_{\{t\}} \Delta(\tilde{P}^s + t e_{s+i}) = \Delta(\tilde{P}^s), \quad i = 1, 2, \dots, r; \quad \gamma_{r+1} < \Delta(\tilde{P}^s).$$

If $\{t_{jk}^{s+i}\}$ do not converge to zero, then, putting $\lim_{l \rightarrow \infty} t_{k_l}^{s+i} = t^{s+i} \neq 0$, we find

$$\lim_{l \rightarrow \infty} \Delta(P_{j k_l}^s + t_{j k_l}^{s+i} e_{s+i}) = \Delta(\tilde{P}^s + t^{s+i} e_{s+i}) = \Delta(\tilde{P}^s),$$

which is impossible. Consequently,

$$\lim_{k \rightarrow \infty} t_{jk}^{s+i} = 0, \quad i = 1, 2, \dots, r.$$

We now have the estimate

$$d = \Delta(\tilde{P}^s) - \gamma_{r+1} \leq \Delta(P_{jk}^s + t_{jk}^{s+1} e_{s+1} + t_{jk}^{s+2} e_{s+2} + \dots + t_{jk}^{s+r} e_{s+r} + t^0 e_{s+r+1}) -$$

$$-\Delta(\tilde{P}^s + t^0 e_{s+r+1}) \leq \text{const} \left[\rho(P_{jk}^s, \tilde{P}^s) + \sum_{i=1}^r |t_{jk}^{s+i}| \right] \rightarrow 0,$$

which is impossible. Hence, at the point \tilde{P}^s there is no decrease in any of the coordinate directions. This means that the tangent plane to the surface Δ at the point $(\tilde{\Delta}, \tilde{P}^s)$, $\tilde{\Delta} = \Delta(\tilde{P}^s)$, has the equation $\Delta = \tilde{\Delta}$. But, by virtue of the

convexity of Δ , the tangent plane is simultaneously a supporting plane, so that all points of Δ , except $(\widetilde{\Delta}, \widetilde{P}^s)$, are situated above the plane $\Delta = \widetilde{\Delta}$. This is equivalent to the assertion of the theorem.

In the case when E_1 and E_2 are Hilbert spaces and $\Delta(P)$ is a function of the form (3), one can prove more precise theorems than Theorems 2 and 3. Namely, it turns out that the corresponding points P_j and P_j^s always converge to quite definite points, depending, generally speaking, on the initial point (see (3)).

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Note: Figure translations are in progress. See original paper for figures.

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