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# MATHEMATICS

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**Abstract**

**Full Text**

**MATHEMATICS**

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## ON THE THEORY OF THE METHODS OF NEWTON AND CHAPLYGIN

*(Presented by Academician S. L. Sobolev, 14 I 1958)*

1. In an analogue of Newton's method <sup>(1)</sup> for the approximate solution of the equation  $P(x) = 0$ , the algorithm

$$x_{n+1} = x_n - \Gamma_n^{-1}P(x_n), \quad (1)$$

is used, where  $\Gamma_x = P'(x_n)$  or  $\Gamma_n = P'(x_0)$ . Instead of the derivative operators  $\Gamma_n$  in (1), one may take an operator  $L$ , having an inverse  $L^{-1}$ , "close" to  $P'(x)$  in the Banach norm <sup>(2)</sup> (our notation). We may impose weaker restrictions on the operator  $L$  and choose a simpler (not necessarily invertible) operator convenient for computations.

Let  $X$  be a space of type  $(B_K)$  <sup>(3)</sup>. We shall use the notation:

$$V^n(z) = V[V^{n-1}(z)], \quad V^0 = I, \quad I(x) \equiv x, \quad \Delta U = U(x + \Delta x) - U(x).$$

**Theorem 1.** Suppose that on the set  $G \subset X$ : 1)  $|\Delta U| \leq V(|\Delta x|)$ , where  $U = I - LP$ ; 2) the operator  $V$  is monotonically increasing; 3)  $G$  contains the bounded sphere  $(x_0, R)$ :

$$|x - x_0| \leq R = \sum_{k=0}^{\infty} V^k(z), \quad z = |LP(x_0)|; \quad (2)$$

- 4) the operator  $LP$  is  $(bk)$ -continuous; 5) the equation  $L(y) = 0$  has the unique zero solution.

Then the algorithm

$$x_{n+1} = x_n - LP(x_n) \quad (n = 0, 1, 2, \dots) \quad (3)$$

converges in the sphere  $(x_0, R)$  to the solution  $x^*$  of the equation  $P(x) = 0$  with rate

$$|x_n - x^*| \leq \sum_{k=n}^{\infty} V^k(z). \quad (4)$$

For condition 2) it is sufficient that the operator  $V$  be additive; for 4), (o)-continuity of  $V$  is sufficient; for 5), additivity and invertibility of the operator  $L$  are sufficient. If  $V$  is not additive, then in the inequalities (2), (4) one may take  $|V_k|$ , and, in general, a stronger estimate is obtained than when using the ordinarily employed (in Banach spaces) power  $\|V\|^k$  or even the norm  $\|V^k\|$  of the iteration.

The operator  $V$  is called **Volterra** if  $(I - V)^{-1} = \sum_{k=0}^{\infty} V^k$ . In this case  $R = (I - V)^{-1}(z)$ . If  $V$  is additive, then  $|x_n - x^*| \leq V^n(R)$ .

2. In what follows  $X$  is a  $K$ -space <sup>(3)</sup>. For practical estimates of a solution it is sometimes necessary to establish an operator analogue of Chaplygin's theorem—

Lygina [4] on differential inequalities, i.e., the establishment of conditions sufficient for the inequality  $P(x) \geq 0$  to imply  $x \geq x^*$ , where  $P(x^*) = 0$ . For an additive invertible operator  $P$  this means its positive invertibility, i.e.,  $P^{-1} > 0$ .

**Theorem 2** (comparison of additive operators). *If the operators  $\Gamma$  and  $\Lambda$  are additive and invertible,  $\Gamma > \Lambda$ ,  $\Gamma^{-1} > 0$ ,  $(I - \Gamma^{-1}\Lambda)^n(x) \xrightarrow{(o)} 0$  ( $n \rightarrow \infty$ ), then  $\Lambda^{-1} > 0$ .*

**Theorem 3** (on operator inequalities). *Suppose that on the set  $G \subset X$  the following conditions hold:  $P(x_0) > 0$  (or  $< 0$ );  $\Gamma(\Delta x) \geq \Delta P \geq \Lambda(\Delta x)$  for  $\Delta x > 0$ ; the operators  $\Gamma$  and  $\Lambda$  are additive and positively invertible; the operator  $\Gamma^{-1}P$  is monotonically continuous; the set  $G$  contains  $[x_1, x_0]$  (or  $[x_0, x_1]$ ), where  $x_1 = x_0 - \Lambda^{-1}P(x_0)$ .*

*Then on this interval there exists a unique solution of the equation  $P(x) = 0$ .*

From this, in particular, it is not difficult to obtain the comparison theorem from the paper [5].

3. Let us construct an algorithm of Chaplygin type on the basis of replacing the operator  $\Gamma_n^{-1}$  in the algorithm of Theorem 1 of the paper [6], similarly to how this was done above.

**Theorem 4.** *Suppose that on each  $[a, b] \subseteq [x_0, \bar{x}_0]$  there exist additive positively invertible operators  $\underline{\Gamma}(a, b)$  and  $\bar{\Gamma}(a, b)$  such that throughout  $[a, b]$*

$$\underline{\Gamma}(a, b)(x - a) \geq P(x) - P(a), \quad \bar{\Gamma}(a, b)(b - x) \geq P(b) - P(x).$$

*Suppose there exists an additive positively invertible operator  $\Gamma \geq \underline{\Gamma}(a, b)$ ,  $\Gamma \geq \bar{\Gamma}(a, b)$  on all  $[a, b]$ , the operator  $\Gamma^{-1}P$  is monotonically continuous, and  $P(x_0) \leq 0 \leq P(\bar{x}_0)$ .*

Define the algorithm

$$\underline{x}_{n+1} = \underline{x}_n - \underline{L}_n(\underline{z}_n), \quad \bar{x}_{n+1} = \bar{x}_n - \bar{L}_n(\bar{z}_n), \quad (5)$$

where  $P(\underline{x}_n) \leq \underline{z}_n \leq 0 \leq \bar{z}_n \leq P(\bar{x}_n)$ , the operators  $L_n$  are positive and homogeneous,

$$\underline{\Gamma}_n \underline{L}_n \leq I, \quad \bar{\Gamma}_n \bar{L}_n \leq I, \quad \underline{\Gamma}_n = \underline{\Gamma}(\underline{x}_n, \bar{x}_n), \quad \bar{\Gamma}_n = \bar{\Gamma}(\underline{x}_n, \bar{x}_n).$$

Then the algorithm determines sequences  $x_n$  satisfying the inequalities

$$\underline{x}_n \leq \underline{x}_{n+1} \leq \underline{x} \leq \bar{x} \leq \bar{x}_{n+1} \leq \bar{x}_n,$$

where  $\underline{x}, \bar{x}$  are the least and greatest solutions on  $[x_0, \bar{x}_0]$  of the equation  $P(x) = 0$ .

For uniqueness of the solution it is sufficient that there exist an additive positively invertible operator  $\Lambda$  such that  $\Delta P \geq \Lambda(\Delta x)$  for  $\Delta x > 0$ .

If it is known in advance that the solution is unique, then the condition  $x_0 \leq \bar{x}_0$  need not be imposed. Instead of the conditions imposed on  $L_n$ , it is sufficient to require that  $\underline{\Gamma}_n \underline{L}_n(\underline{z}_n) \geq \underline{z}_n, \bar{\Gamma}_n \bar{L}_n(\bar{z}_n) \leq \bar{z}_n$ . Under some additional assumptions concerning  $L_n(z_n)$ , the algorithm converges to the solution. Instead of continuity of the operator  $\Gamma^{-1}P$ , one may require continuity of the operators  $\Gamma$  and  $P$ . This remark also applies to Theorem 3.

Putting  $\Gamma_n = \bar{\Gamma}_n, L_n = \Gamma_n^{-1}$ , we obtain the algorithms of the paper [7].

**Theorem 5.** *If, under the conditions of Theorem 4, the operators  $I - \Gamma_n$  are positive Volterra operators, then one may take as  $L_n$*

$$\sum_{k=0}^m (I - \Gamma_n)^k$$

where  $m$  is arbitrary.

Hence, for  $z_n = P(x_n)$ , acting in the space of collections of derivatives, one can, in particular, obtain (concrete) algorithms of work<sup>8</sup>. They occupy an intermediate place between the algorithms

$$x_{n+1} = x_n - P(x_n)$$

and the more complicated, but more rapidly convergent,

$$x_{n+1} = x_n - \Gamma_n^{-1}P(x_n).$$

4. The additivity and positive invertibility condition on the operator  $\Gamma'$ , usually used in analogues of Chaplygin's method<sup>6,7</sup>, can be replaced by another one.

**Theorem 6.** *Let the following conditions be fulfilled on the set  $G \subset X$ :  $P(x_0) > 0$  (or  $< 0$ ),  $\Gamma(\Delta x) \geq \Delta P$  for  $\Delta x > 0$ ; there exists an operator  $L$  such that  $L > 0$ ,  $\Gamma L \leq I$  (or  $L(y) < 0 \leq \Gamma L(y) - y$  for  $y < 0$ ); the equation  $L(y) = 0$  has the unique zero solution; the operator  $LP$  is monotonically continuous;  $\Delta U \leq \Delta V$  for  $\Delta x > 0$ , where  $U = I - LP$ ; the operator  $V$  is monotonically increasing; the set  $G$  contains the bounded interval*

$$[x_0 - R, x_0] \quad (\text{or } [x_0, x_0 + R]),$$

where

$$R = \sum_{k=0}^{\infty} V^k(z), \quad z = |LP(x_0)|.$$

Then the equation  $P(x) = 0$  has on this interval a solution  $x^*$ , which can be obtained by algorithm (3), and moreover  $x_n \searrow x^*$  (or  $x_n \nearrow x^*$ ) with rate

$$|x_n - x^*| \leq \sum_{k=n}^{\infty} V^k(z).$$

A consequence of this is Theorem 3 under the additional condition: the operator  $V = I - \Gamma^{-1}\Lambda$  is Volterra. The theorem can be modified by replacing algorithm (3) with an algorithm of the form (5).

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*Note: Figure translations are in progress. See original paper for figures.*

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